

# S-Packing Colorings of Cubic Graphs

Nicolas Gastineau, Olivier Togni

#### ▶ To cite this version:

Nicolas Gastineau, Olivier Togni. S-Packing Colorings of Cubic Graphs. 2014. hal-00967446v1

# HAL Id: hal-00967446 https://u-bourgogne.hal.science/hal-00967446v1

Preprint submitted on 28 Mar 2014 (v1), last revised 29 Apr 2016 (v2)

**HAL** is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers. L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

# S-Packing Colorings of Cubic Graphs

Nicolas Gastineau\*and Olivier Togni

LE2I, UMR CNRS 6303 Université de Bourgogne, 21078 Dijon cedex, France March 28, 2014

Given a non-decreasing sequence  $S=(s_1,s_2,\ldots,s_k)$  of positive integers, an S-packing coloring of a graph G is a mapping c from V(G) to  $\{s_1,s_2,\ldots,s_k\}$  such that any two vertices with color  $s_i$  are at mutual distance greater than  $s_i, 1 \leq i \leq k$ . This paper studies S-packing colorings of (sub)cubic graphs. We prove that subcubic graphs are (1,2,2,2,2,2,2)-packing colorable and (1,1,2,2,3)-packing colorable. For subdivisions of subcubic graphs we derive sharper bounds, and we provide an example of a cubic graph of order 38 which is not  $(1,2,\ldots,12)$ -packing colorable.

**Keywords:** coloring, packing chromatic number, cubic graph.

#### 1 Introduction

A proper coloring of a graph G is a mapping which associates a color (integer) to each vertex such that adjacent vertices get distinct colors. In such a coloring, the color classes are stable sets (1-packings). As an extension, a d-distance coloring of G is a proper coloring of the d-th power  $G^d$  of G, i.e. a partition of V(G) into d-packings (sets of vertices at pairwise distance greater than d). While Brook's theorem implies that all cubic graphs except the complete graph  $K_4$  of order 4 are properly 3-colorable, many authors studied 2-distance colorings of cubic graphs.

The aim of this paper is to study a mixing of these two types of colorings, i.e. colorings of (sub)cubic graphs in which some colors classes are 1-packings while other are d-packings,  $d \geq 2$ . Such colorings can be expressed using the notion of S-packing coloring: For a non-decreasing sequence  $S = (s_1, s_2, \ldots, s_k)$  of positive integers, an S-packing coloring (or simply S-coloring) of a graph G is a coloring of its vertices with colors from

<sup>\*</sup>Author partially supported by the Burgundy Council

 $\{s_1, s_2, \ldots, s_k\}$  such that any two vertices with color  $s_i$  are at mutual distance greater than  $s_i$ ,  $1 \le i \le k$ . The color class of each color  $s_i$  is thus an  $s_i$ -packing. The graph G is S-colorable if there exists an S-coloring and it is S-chromatic if it is S-colorable but not S'-colorable for any  $S' = (s_1, s_2, \ldots, s_j)$  with j < k (note that Goddard et al. [13] define differently the S-chromaticness for infinite graphs).

A  $(d, \ldots, d)$ -coloring is thus a d-distance k-coloring, where k is the number of d (see [16] for a survey of results on this invariant) while a  $(1, 2, \ldots, d)$ -coloring is a packing coloring. The packing chromatic number  $\chi_{\rho}(G)$  of G is the integer k for which G is  $(1, \ldots, k)$ -chromatic. This parameter was introduced recently by Goddard et al. [11] under the name of broadcast chromatic number and the authors showed that deciding whether  $\chi_{\rho}(G) \leq 4$  is NP-hard. A series of works [3, 5, 7, 8, 11, 17] considered the packing chromatic number of infinite grids. For sequences S other than  $(1, 2, \ldots, k)$ , S-packing colorings were considered more recently [6, 10, 12, 13].

Regarding cubic graphs, the packing chromatic number of the hexagonal lattice and of the infinite 3-regular tree is 7 and at most 7, respectively. Goddard et al. [11] asked what is the maximum of the packing chromatic number of a cubic graph of order n. For 2-distance coloring of cubic graphs, Cranston and Kim have recently shown [4] that any subcubic graph is (2, 2, 2, 2, 2, 2, 2, 2, 2)-colorable (they in fact proved a stronger statement for list coloring). For planar subcubic graphs G, there are also sharper results depending on the girth of G [2, 4, 15].

In this paper, we study S-packing colorings of subcubic graphs for various sequences S starting with one or two '1'. We also compute the distribution of S-chromatic cubic graphs up to 20 vertices, for three sequences S. The corresponding results are reported on Tables 1, 2, and 3. They are obtained by an exhaustive search, using the lists of cubic graphs maintained by Gordon Royle [14]. The paper is organized as follows: Section 2 is devoted to the study of  $(1, k, \ldots, k)$ -colorings of subcubic graphs for k = 2 or 3; Section 3 to  $(1, 1, 2, \ldots)$ -colorings; Section 4 to  $(1, 2, 3, \ldots)$ -colorings and Section 5 concludes the paper by listing some open problems.

### 1.1 Notation

To describe an S-coloring, if an integer s is repeated in the sequence S, then we will denote the colors s by  $s_a, s_b, \ldots$ 

The subdivided graph S(G) of a (multi)graph G is the graph obtained from G by subdividing each edge once, i.e. replacing each edge by a path of length two. In S(G), vertices of G are called original vertices and other vertices are called subdivision vertices. Let us call a graph d-irregular if it has no adjacent vertices of degree d. Note that graphs obtained from subcubic graphs by subdividing each edge at least once are 3-irregular graphs.

The following method (that is inspired from that of Cranston and Kim [4]) is used in the remainder of the paper to produce a desired coloring of a subcubic graph: for a graph G and an edge  $e = xy \in E(G)$ , a level ordering of (G, e) is a partition of V(G) into levels  $L_i = \{v \in V(G) : d(v, e) = i\}, 0 \le i \le \epsilon(e)$ , with  $\epsilon(e) = \max(\{d(u, e), u \in V(G)\}) \le diam(G)$ . The vertices are then colored one by one, from level  $\epsilon(e)$  to 1, while

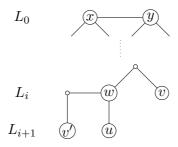


Figure 1: Level ordering of a subcubic graph from the edge xy. The vertices v, w are siblings and v and v' are both cousins of u.

preserving some properties. These properties are used at the end to allow to color the vertices x and y by recoloring possibly some vertices in their neighborhood.

Two vertices u and v of G are called *siblings* if they are not adjacent, are on the same level  $L_i$  for some  $i \geq 1$  and have a common neighbor in  $L_{i-1}$ . Two vertices u and v of G are called *cousins* if they are at distance 3 and in every path of length 3 between u and v, there is a neighbor of u or v in a lower level than the level of u and v. Note that a vertex has at most one sibling and two cousins (see Figure 1). Given a (partial) coloring c of G, let  $C_1(u) = \{c(v) : uv \in E(G)\}$ ,  $C_2(u) = \{c(v) : d(u,v) = 2$ , with u,v not siblings $\{cunu}$ , with  $\{cunu}$ , with  $\{cunu}$  and  $\{cunu}$ , with  $\{cunu}$ , w

## 2 $(1, k, \ldots, k)$ -coloring

In this section,  $(1, k, \dots, k)$ -colorings of subcubic graphs are studied for k = 2 or 3.

#### **2.1** $(1,3,\ldots,3)$ -coloring

The following proposition is used to obtain an S-coloring of a subdivided graph:

**Proposition 1.** Let G be a graph and  $S = (s_1, ..., s_k)$  be a non-decreasing sequence of integers. If G is S-colorable then S(G) is  $(1, 2s_1 + 1, ..., 2s_k + 1)$ -colorable.

Proof. Let c be an S-coloring of G. Every pair of vertices  $u, v \in V(G)$  such that d(u, v) = d become at distance 2d + 1 in S(G). Therefore, every set of vertices in V(G) forming an i-packing also forms a 2i + 1-packing in S(G). Using color 1 on subdivision vertices and using the coloring c (considering the sequence differently) on original vertices, we obtain a  $(1, 2s_1 + 1, \ldots, 2s_k + 1)$ -coloring of S(G).

Corollary 1. For every subcubic graph G, S(G) is (1,3,3,3)-colorable.

*Proof.* Brooks' theorem asserts that every subcubic graph except  $K_4$  is (1, 1, 1)-colorable. Hence, by Proposition 1, every subcubic graph G except  $K_4$  is such that S(G) is

(1,3,3,3)-colorable. We define a (1,3,3,3)-coloring of  $S(K_4)$  as follows: let  $\gamma: E(K_4) \to \{a,b,c\}$  be a proper edge 3-coloring of  $K_4$ . Put color 1 on all four original vertices of  $K_4$  and put color  $3_{\gamma(e)}$  on each subdivision vertex corresponding with edge e of  $K_4$ .

Goddard et al. [11] characterized (1,3,3)-colorable graphs as the graphs obtained from any bipartite multigraph by subdividing it and adding leaves on original vertices. Therefore, there are many subdivided subcubic graphs that are not (1,3,3)-colorable (for instance  $S(C_3) = C_6$ ), showing that the bound of Corollary 1 is tight in a certain sense.

#### **2.2** $(1, 2, \dots, 2)$ -coloring

Note that no cubic graph with more than 3 vertices is (1,2,2)-colorable since a graph with three vertices of degree larger than 2 at mutual distance less than 2, is not (1,2,2)-colorable. However, there exist (1,2,2)-colorable subcubic graphs and it has been recently proved [9] that determining if a subcubic bipartite graph is (1,2,2)-colorable is NP-complete.

**Proposition 2.** Every subcubic graph is (1, 2, 2, 2, 2, 2, 2, 2)-colorable.

*Proof.* Let G be a subcubic graph and let e = xy be any edge of G. Define a level ordering  $L_i$ ,  $0 \le i \le r = \epsilon(e)$ , of (G, e).

We first construct a coloring c of the vertices of G from level r to 1 and with colors from the set  $C = \{1, 2_a, 2_b, 2_c, 2_d, 2_e, 2_f\}$ , that satisfies the following properties:

- i) color 1 is used as often as possible, i.e. when coloring a vertex u, if no neighbor is colored 1, then u is colored 1;
- ii) if u is colored 2, then there is subsidiary color  $\tilde{c}(u) \in C$  different from c(u) such that  $\tilde{c}(u) \notin C_1(u) \cup C_2(u)$ , but with possibly  $\tilde{c}(u) = c(v)$  if u and v are siblings.

The set  $L_r$  induces a disjoint union of paths and cycles in G. Since paths and cycles are (1,2,2,2)-colorable, we are able to construct a coloring of the vertices of  $L_r$  as follows. Start by coloring each path/cycle with colors  $\{1,2_a,2_b,2_c\}$ . For each pair of vertices u,v in different paths/cycles at distance 2 both colored  $2_a$  ( $2_b$  or  $2_c$  respectively), set  $c(u) = 2_d$  ( $2_e$  or  $2_f$ , respectively). Afterwards, for every vertex u of colors  $2_a$  ( $2_b$ ,  $2_c$ ,  $2_d$ ,  $2_e$  or  $2_f$ , respectively), set  $\tilde{c}(u) = 2_d$  ( $2_e$ ,  $2_f$ ,  $2_a$ ,  $2_b$  or  $2_c$ , respectively). Then Property ii) is satisfied.

Assume that we have already colored all vertices of G of levels from r to i+1 and that we are going to color vertex  $u \in L_i$ ,  $1 \le i \le r-1$ . If  $1 \notin C_1(u)$  then set c(u) = 1 (Property i) is then satisfied). Now, if  $1 \in C_1(u)$ , then let  $u_1$  be the neighbor of u of color 1 and let  $u_2$  be the other neighbor of u, if any. By construction, either  $c(u_2) = 1$  or  $1 \in C_1(u_2)$ , hence  $|C_1(u) \cup C_2(u)| \le 5$ . In that case there are at least two colors  $\{2_{\alpha}, 2_{\beta}\} \subset C \setminus C_2(u)$  for some  $\alpha, \beta \in \{a, \ldots, f\}$ , with possibly, if u has a sibling v,  $2_{\beta} = c(v)$ . Then set  $c(u) = 2_{\alpha}$  and  $\tilde{c}(u) = 2_{\beta}$  (Property ii) is then satisfied). Figure 2 illustrates this case.

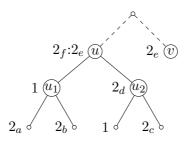


Figure 2: A configuration in the proof of Proposition 2, when coloring vertex u. The label  $2_f:2_e$  on u means that  $c(u)=2_f$  and  $\tilde{c}(u)=2_e$ .

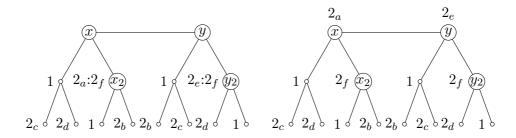


Figure 3: A configuration in the proof of Proposition 2, before (on the left) and after (on the right) coloring x and y.

Finally, it remains to color vertices of  $L_0$ , i.e., x and y. If  $1 \in C_1(x) \cap C_1(y)$  then, by Property i), the neighbor  $x_2$  of x colored 2, if any, has a neighbor of color 1 and the same goes for y, with  $y_2$  being the neighbor of y colored 2, if any. Hence  $|C_1(x) \cup C_2(x)| \leq 5$ ,  $|C_1(y) \cup C_2(y)| \leq 5$  and there remains at least a color  $2_{\alpha}$  available for x and a color  $2_{\beta}$  for y. If  $\alpha = \beta$  then set  $c(x) = c(x_2)$ ,  $c(y) = c(y_2)$ ,  $c(x_2) = \tilde{c}(x_2)$  and  $c(y) = \tilde{c}(y_2)$ . Figure 3 illustrates that case.

If  $1 \in C_1(x)$  but  $1 \notin C_1(y)$  (or  $1 \in C_1(y)$  but  $1 \notin C_1(x)$ , by symmetry), then set c(y) = 1 and if  $C_2(x) = C$  then set  $c(x) = c(x_2)$  and  $c(x_2) = \tilde{c}(x_2)$ , else give to x an available color.

Otherwise,  $1 \notin C_1(x) \cup C_1(y)$ . Then set c(y) = 1 and we show that there is always a color 2 to assign to x. If  $|C_1(x) \cup C_2(x)| \le 6$  then there is a color available for x. Else, let  $x_1, x_2$  be the two neighbors of x other than y and let  $x_1'$  ( $x_2'$ , respectively) be the neighbor of  $x_1$  ( $x_2$ , respectively) colored 2 other than x (no more than one, as  $x_1$  and  $x_2$  both have a neighbor colored 1). Suppose, without loss of generality, that  $c(x_1) = 2_a$ ,  $c(x_2) = 2_b$ ,  $c(x_1') = 2_c$  and  $c(x_2') = 2_d$ . If  $\tilde{c}(x_1) \in \{2_d, 2_e, 2_f\}$  then recolor  $x_1$  by its subsidiary color  $\tilde{c}(x_1)$  and set  $c(x) = 2_a$ . Similarly, if  $\tilde{c}(x_2) \in \{2_c, 2_e, 2_f\}$  then recolor  $x_2$  by its subsidiary color  $\tilde{c}(x_2)$  and set  $c(x) = 2_b$ . Else,  $\tilde{c}(x_1) = 2_b$  and  $\tilde{c}(x_2) = 2_a$ . Recolor  $x_1'$  by its subsidiary color  $\tilde{c}(x_1')$  and set  $c(x) = 2_c$ . If  $\tilde{c}(x_1') = 2_a = c(x_1)$ , then switch the colors of  $x_1$  and  $x_2$  (this is possible since  $\tilde{c}(x_1) = c(x_2)$  and  $\tilde{c}(x_2) = c(x_1)$ ). Figure 4

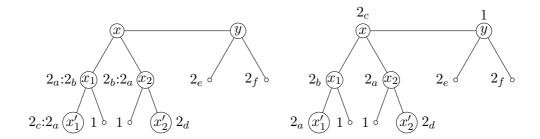


Figure 4: A configuration in the proof of Proposition 2, before (on the left) and after (on the right) coloring x and y.

illustrates this case.

Therefore, we obtain, in all cases, a (1,2,2,2,2,2,2)-coloring of G.

The Petersen graph is an example of cubic graph which is not (1, 2, 2, 2, 2, 2)-colorable, showing that the result of Proposition 2 is tight in a certain sense. However, the experiments reported on Table 1 suggest that the Petersen graph could be the only non (1, 2, 2, 2, 2, 2)-colorable subcubic graph.

$n \backslash S$	(1, 2, 2, 2)	(1, 2, 2, 2, 2)	(1, 2, 2, 2, 2, 2)	(1, 2, 2, 2, 2, 2, 2)
4	1	0	0	0
6	1	1	0	0
8	2	1	2	0
10	11	7	0	1
12	11	74	0	0
14	254	250	5	0
16	1031	3017	12	0
18	15960	25297	44	0
20	178193	332045	251	0
22	2481669	4835964	1814	0

Table 1: Number of S-chromatic cubic graphs of order n up to 22.

Furthermore, as the following proposition shows, even some bipartite cubic graphs are not (1, 2, 2, 2, 2, 3)-colorable.

**Proposition 3.** There exist bipartite cubic graphs that are not (1, 2, 2, 2, 2, 3)-colorable.

*Proof.* The cubic graph depicted in Figure 5 is bipartite and is (1, 2, 2, 2, 2, 2)-colorable, as shown on the figure. Let (A, B) be the two subsets of vertices that form a bipartition of this graph. Suppose this graph is (1, 2, 2, 2, 2, 3)-colorable and let c be a (1, 2, 2, 2, 2, 3)-coloring and  $X_1$  be the set of vertices colored 1. Remark that the cardinality of any 2-packing is at most 2 and that any pair of vertices (u, v) included in A or in B is such

that  $d(u,v) \leq 2$ . We have  $|X_1| \geq 5$ , as at most one vertex can be colored 3 (since the diameter of the graph is 3) and at most two vertices can be colored the same color 2.

Firstly, if  $X_1 \subseteq A$  or  $X_1 \subseteq B$ , then each remaining vertex should be colored differently in the other partition, which is impossible since |A| = |B| = 7.

Secondly, if there are vertices colored 1 in A and B, then the only possibility in order to have  $|X_1| \geq 5$  is to have one vertex colored 1 in one partition and four vertices colored 1 in the other partition. Suppose, without loss of generality, that  $|X_1 \cap A| = 1$  and  $|X_1 \cap B| = 4$ . Exactly three vertices are not colored 1 in B. Consequently, only three pairs of vertices can have the same color 2 and the nine vertices not colored 1 cannot be all colored with the remaining colors 2 and 3.

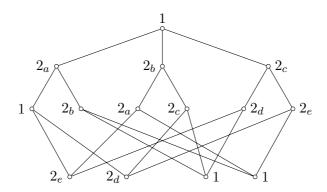


Figure 5: A cubic bipartite (1, 2, 2, 2, 2, 2)-chromatic graph of order 14.

The next results show that there are sub-families of subcubic graphs that can be colored with fewer colors.

**Proposition 4.** Every 3-irregular subcubic graph is (1, 2, 2, 2)-colorable.

*Proof.* Let G be a 3-irregular graph and let e = xy be any edge of G such that x and y are both of degree at most 2. If no such edge exists, then the graph is the subdivision S(H) of some subcubic graph H where leaves could be added on original vertices of degree 2 and thus G is (1,3,3,3)-colorable by Corollary 1. Define a level ordering  $L_i$ ,  $0 \le i \le r = \epsilon(e)$ , of (G,e).

We construct a coloring c of the vertices of G from level r to 1 and with colors from the set  $\{1, 2_a, 2_b, 2_c\}$ , that satisfies the following properties:

- i) color 1 is used as often as possible, i.e. when coloring a vertex u, if no neighbor is colored 1, then u is colored 1;
- ii) every vertex of degree 2 is colored 1 when first coloring vertices of  $L_i$ , except if the connected component containing this vertex in  $L_i$  is a path of order 2 (in which case one of the two vertices is colored 1).

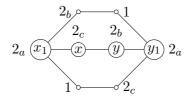


Figure 6: The graph  $G_{(1,2,2,2)}$  from Proposition 4.

The set  $L_r$  induces a disjoint union of paths of order at most 3 in G. Since paths are (1,2,2)-colorable,  $L_r$  is (1,2,2)-colorable. Moreover, in every path of order 3 in  $L_r$ , the central vertex has degree 3, thus a color 1 could be given to every vertex of degree 2. If the path is of order 2, one of the vertex is colored 1. Thus, Properties i) and ii) are satisfied.

Assume that we have already colored all vertices of G of levels from r to i+1 and that we are going to color vertex  $u \in L_i$ ,  $1 \le i \le r-1$ . We consider two cases depending on the degree of u:

#### Case 1. u is of degree 3.

If  $1 \notin C_1(u)$ , then u can be colored 1. Let  $u_1$  and  $u_2$  be the colored neighbors of u, with  $c(u_1) = 1$ . By Property i), a colored neighbor of  $u_2$  has color 1. Hence, we have  $|C_1(u) \cup C_2(u)| \leq 3$  and u can be colored some color 2.

#### Case 2. u is of degree at most 2.

Let  $u_1$  be the colored neighbor of u, if any. If  $u_1$  is of degree 3, let  $u_{1,1}$  and  $u_{1,2}$  be the colored neighbor of u,  $u_{1,1,1}$  be the neighbor of  $u_{1,1}$  different from u and  $u_{1,2,1}$  be the neighbor of  $u_{1,2}$  different from u. If  $1 \notin C_1(u)$ , then we can set c(u) = 1. Otherwise,  $c(u_1) = 1$  and thus  $c(u_{1,1}) \neq 1$  and  $c(u_{1,2}) \neq 1$ . Therefore,  $c(u_{1,1,1}) = c(u_{1,2,1}) = 1$ . Thus,  $u_1$  can be recolored some color 2 and we can set c(u) = 1. If  $u_1$  is of degree at most 2, then, as  $|C_1(u_1) \cup C_2(u_1)| \leq 3$ , we can recolor  $u_1$  by a color 2. Thus, we can set c(u) = 1.

Finally, it remains to color vertices of  $L_0$ , i.e. x and y. Let  $x_1$  be the possible neighbor of x different from y and let  $y_1$  be the possible neighbor of y different from x. We consider three cases that cover all the possibilities by symmetry:

#### Case 1. $x_1$ and $y_1$ both have degree 3.

If their neighbors different from x and y are not adjacent between them, then, by Property ii), these vertices have color 1 and  $x_1$  and  $y_1$  have some color 2. Thus we can set c(x) = 1 and some color 2 to y, as  $|C_1(y) \cup C_2(y)| \leq 3$ . Suppose two pairs of neighbors of  $x_1$  and  $y_1$  different from x and y are adjacent. Thus, this graph is the graph  $G_{(1,2,2,2)}$  from Figure 6 and is (1,2,2,2)-colorable. Suppose only two neighbors of  $x_1$  and  $y_1$  different from x and y are adjacent. Let  $x_{1,1}$  and  $y_{1,1}$  be these two neighbors, the other neighbors are colored 1 by Property ii). One of

these two vertices is colored 1 and the other one is colored 2. Suppose without loss of generality that  $c(x_{1,1}) = 1$ . Hence, we have  $|C_1(x) \cup C_2(x)| \leq 3$  and we can color x by a color 2 and set c(y) = 1.

**Case 2.**  $x_1$  has degree at most 2 and  $y_1$  has degree 3.

By Property ii),  $x_1$  is colored 1. As  $|C_1(y_1) \cup C_2(y_1)| \leq 3$ , then  $y_1$  can be recolored some color 2. Thus, we can set c(y) = 1, and as  $|C_1(x) \cup C_2(x)| \leq 3$  we can set a color 2 to x.

Case 3.  $x_1$  and  $y_1$  are both of degree at most 2.

If  $x_1$  and  $y_1$  are adjacent, then the graph is  $C_4$  which is trivially (1, 2, 2, 2)-colorable. If  $x_1$  and  $y_1$  are not adjacent, then they both have color 1,  $|C_1(x) \cup C_2(x)| \leq 2$  and  $|C_1(y) \cup C_2(y)| \leq 2$ . Thus, we can set some colors 2 to x and y.

Therefore, we obtain in all cases a (1, 2, 2, 2)-coloring of G.

Remark that the 5-cycle  $C_5$  is 3-irregular and is not (1,2,2)-colorable, hence the result of Proposition 4 is tight in a certain sense. However, there are 3-irregular subcubic graphs that are (1,2,2,3)-colorable. The graph from Figure 6 is such an example (the vertex x can be recolored 1 and then color  $2_c$  can be replaced by color 3).

We end this section with some results on subdivided graphs. Let  $\delta(G)$  be the minimum degree of G.

**Proposition 5.** For every graph G with  $\delta(G) \geq 3$ , if S(G) is (1,2,2)-colorable then G is bipartite.

*Proof.* Suppose S(G) is (1,2,2)-colorable and G contains an odd cycle. In every (1,2,2)-coloring of a graph, every vertex of degree at least 3 should be colored some color 2 (if a vertex of degree at least 3 is colored 1, the coloring cannot be extended to the neighbors of this vertex). Therefore, if G contains an odd cycle, then S(G) contains a cycle with an odd number of vertices of degree 3 and the colors  $2_a$  and  $2_b$  are not sufficient to alternately color these vertices. Hence S(G) is not (1,2,2)-colorable.

As every bipartite graph G is (1,1)-colorable, then by Proposition 1, S(G) is (1,3,3)-colorable (and also (1,2,2)-colorable and (1,2,3)-colorable). Thus, we obtain the following corollary.

Corollary 2. For every graph G with  $\delta(G) \geq 3$ ,

S(G) (1,2,2)-colorable  $\Leftrightarrow S(G)$  (1,2,3)-colorable  $\Leftrightarrow S(G)$  (1,3,3)-colorable  $\Leftrightarrow G$  bipartite.

# **3** $(1, 1, 2, \ldots)$ -coloring

Remind that bipartite graphs are (1,1)-colorable. For non-bipartite subcubic graphs, we prove the following:

**Proposition 6.** Every subcubic graph is (1, 1, 2, 2, 3)-colorable.

*Proof.* Let G be a subcubic graph and let e = xy be any edge of G. Define a level ordering  $L_i$ ,  $0 \le i \le r = \epsilon(e)$ , of (G, e).

We first construct a coloring c of the vertices of G from level r to 1 and with colors from the set  $\{1_a, 1_b, 2_a, 2_b, 3\}$ , that satisfies the following properties:

- i) colors 1 are used as often as possible, i.e. when coloring a vertex u, if no neighbor is colored  $1_a$  then u is colored  $1_a$  else if no neighbor is colored  $1_b$ , then u is colored  $1_b$ ;
- ii) if u is colored 2 or 3, then, except in the case where u and a sibling of u are both colored some color 2, u has a subsidiary color  $\tilde{c}(u) \in \{2_a, 2_b, 3\}$  different from c(u) such that:
  - $\tilde{c}(u) \notin C_1(u) \cup C_2(u) \cup C_3(u)$ , but with possibly  $\tilde{c}(u) \in \tilde{C}_3(v)$ , if  $\tilde{c}(u) = 3$ ;
  - $\tilde{c}(u) \notin C_1(u) \cup C_2(u)$ , otherwise;
- iii) if u is colored 3 then its sibling (if any) is colored  $1_a$  or  $1_b$ .

We begin with recalling that by Property i), vertices of color among  $\{2_a, 2_b, 3\}$  have neighbors at the same level or at the above level colored  $1_a$  and  $1_b$ . Also, by properties ii) and iii), we have the following claim.

Claim 1. If two cousins u and v are such that  $\tilde{c}(u) = 3$  and c(v) = 3 and there is no vertex at distance at most 2 from v colored by  $\{\tilde{c}(v)\} \cap \{2_a, 2_b\}$  in the below levels and no vertex at distance at most 3 from u colored by 3 in the below levels, then u and v can be recolored in order that c(u) = 3.

The set  $L_r$  induces a disjoint union of paths and cycles in G. Since paths and cycles are (1,1,2)-colorable, we are able to construct a coloring of the vertices of  $L_r$  as follows. Start by coloring each path/cycle with colors  $\{1_a,1_b,2_a\}$ , using one color  $2_a$  per odd cycle (and no color 2 for even cycle). For each pair of vertices in different paths/cycles at distance 2 both colored  $2_a$ , recolor one of the vertex with color  $2_b$ . Then, Properties i), ii) and iii) are satisfied.

Assume that we have already colored all vertices of G of levels from r to i+1 and that we are going to color vertex  $u \in L_i$ ,  $1 \le i \le r-1$ . If  $C_1(u) \ne \{1_a, 1_b\}$ , then give to u the available color 1. Hence Property i) is satisfied and we can suppose that u has two neighbors in  $L_i \cup L_{i+1}$ , say  $u_1$  and  $u_2$ , such that  $c(u_1) = 1_a$  and  $c(u_2) = 1_b$ . Moreover, we can suppose that  $u_1$  has a neighbor  $u_{1,1}$  of color  $1_b$  and  $u_2$  has a neighbor  $u_{2,1}$  of color  $1_a$  since otherwise, u could be colored by either  $1_a$  or  $1_b$  after recoloring  $u_1$  or  $u_2$ . Let  $u_{1,2}$  be the other neighbor of  $u_1$  different from u, if any, with  $c(u_{1,2}) = \alpha$  and let  $u_{2,2}$  be the other neighbor of  $u_2$  different from u, if any, with  $c(u_{2,2}) = \beta$ . We consider three cases depending on the values of  $\alpha$  and  $\beta$ :

**Case 1.**  $\alpha = 1_b$  and  $\beta = 1_a$ . Then set  $c(u) = 2_a$  or  $c(u) = 2_b$  if the sibling of u is colored  $2_a$ , and  $\tilde{c}(u) = 2_b$  if u has no sibling colored 2. Thus Property ii) is satisfied.

- **Case 2.**  $\alpha = 1_b$  and  $\beta \neq 1_a$  (or  $\alpha \neq 1_b$  and  $\beta = 1_a$ , by symmetry). As  $u_{2,2}$  has no sibling colored some color 2, then it has a subsidiary color by Property ii).
  - If  $c(u_{2,2}) \cup \tilde{c}(u_{2,2}) \neq \{2_a, 2_b\}$ , then set  $c(u_{2,2}) = 3$  (possibly, if  $u_{2,2}$  has a cousin z colored 3, it can be seen that  $u_{2,2}$  and z satisfy the conditions of Claim 1, hence z can be recolored by its subsidiary color). Thus we can set  $c(u) = 2_a$  and  $\tilde{c}(u) = 2_b$ , or  $c(u) = 2_b$  if the sibling of u is colored  $2_a$ , and Property ii) is satisfied.
  - If  $c(u_{2,2}) \cup \tilde{c}(u_{2,2}) = \{2_a, 2_b\}$  and the sibling of u is colored 2, then give an appropriate color to  $u_{2,2}$  in order u to have a color different from its sibling.
  - If  $c(u_{2,2}) \cup \tilde{c}(u_{2,2}) = \{2_a, 2_b\}$  and u has no sibling colored 2, suppose, without loss of generality, that  $c(u_{2,2}) = 2_a$ . As  $u_{2,1}$  is colored  $1_a$ , it has a neighbor colored  $1_b$  different from  $u_2$  (if not,  $u_{2,1}$  and  $u_2$  can be recolored and we can give a color 1 to u). If the other neighbor of  $u_{2,1}$  is colored 3 then change its color by its subsidiary color using Claim 1. Since the vertices  $u_{1,1}$  and  $u_{1,2}$  are both colored  $1_b$ , they could have two neighbors colored  $2_a$  and  $2_b$  or these vertices could have one neighbor colored  $1_a$  different from  $u_1$  and another neighbor different from  $u_1$  of any color. In the latter case, recolor the possible neighbor of these vertices colored 3 by its subsidiary color using Claim 1. Then, set  $c(u) = 2_b$  and  $\tilde{c}(u) = 3$ .
- Case 3.  $\alpha, \beta \in \{2_a, 2_b, 3\}$ . Then the set  $A = (c(u_{1,2}) \cup \tilde{c}(u_{1,2})) \cap (c(u_{2,2}) \cup \tilde{c}(u_{2,2}))$  is not empty. If  $3 \in A$ , then set  $c(u_{1,2}) = 3$  and  $c(u_{2,2}) = 3$  and set  $c(u) = 2_a$  and  $\tilde{c}(u) = 2_b$ , or  $c(u) = 2_b$ , if the sibling of u is colored  $2_a$ . Thus Property ii) is satisfied. Else, if  $3 \notin A$  and u has a sibling colored 2, then change the color of  $u_{1,2}$  and  $u_{2,2}$  by appropriate colors and give a color 2 to u. Else, if  $3 \notin A$  and u has no sibling colored 2, then give a color  $2_\delta \in A$ , with  $\delta \in \{a,b\}$  to  $u_{1,2}$  and  $u_{2,2}$  and recolor each vertex at distance at most 3 from u by its subsidiary color using Claim 1 (there are at most two vertices colored 3 by hypothesis). Hence, set  $c(u) = 2_\beta$  and  $\tilde{c}(u) = 3$ ,  $\beta \neq \alpha$  and  $\beta \in \{a,b\}$ , and Property ii) is satisfied. Property iii) is satisfied, as the color 3 has been given to vertices which have no sibling colored 2.

Now, it remains to color vertices of  $L_0$ , i.e., x and y. Let  $x_1$  and  $x_2$  be the possible neighbors of x different from y and  $y_1$  and  $y_2$  be the possible neighbors of y different from x. We consider seven cases that cover all the possible configurations for the colors of the neighbors of x and y (in order to simplify, configurations that can be obtained by exchanging x and y are omitted):

- **Case 1.**  $1_a \notin C_1(x)$  and  $1_b \notin C_1(y)$ . Then set  $c(x) = 1_a$  and  $c(y) = 1_b$ .
- Case 2.  $C_1(x) = \{1_a, \alpha\}$  and  $C_1(y) = \{1_a\}$ , with  $\alpha \in \{2_a, 2_b, 3\}$ . Then set  $c(y) = 1_b$ . Suppose  $c(x_1) = 1_a$  and  $c(x_2) = \alpha$ . The vertex  $x_1$  has a neighbor colored  $1_b$  (if not we would be in Case 1 by recoloring  $x_1$ ). Let  $x_{1,1}$  be the possible neighbor of  $x_1$  not colored  $1_b$  and  $\beta$  be its color. Recolor  $x_{1,1}$  by its subsidiary color if  $\alpha \neq \beta$  and  $\alpha, \beta \in \{2_a, 2_b\}$ . Then give a remaining color 2 to x.

- Case 3.  $C_1(x) = \{\alpha, \beta\}$  and  $C_1(y) = \{1_a\}$ , with  $\alpha, \beta \in \{1_a, 1_b\}$ . Then set  $c(y) = 1_b$ . A vertex among  $x_1$  and  $x_2$  has a neighbor of color  $1_b$  (if not we would be in the first case by recoloring  $x_1$  and  $x_2$ ) and we suppose that this vertex is  $x_1$ . Two cases can occur,  $C_1(x_2) = \{2_a, 2_b\}$  and consequently  $C_2(x_2) = \{1_a, 1_b, 2_a, 2_b\}$  or  $C_1(x_2) = \{\gamma, \delta\}$ , with  $\gamma \in \{1_a, 1_b\}$  and  $\delta \in \{1_a, 1_b, 2_a, 2_b, 3\}$ . In both cases, recolor every vertex at distance at most 3 from x colored 3 by its subsidiary color (by Property iii), such a vertex has no sibling colored 2) and set c(x) = 3.
- **Case 4.**  $C_1(x) = \{1_a, 1_b\}$  and  $C_1(y) = \{2_a, 2_b\}$ . Then set  $c(y) = 1_b$ . We suppose, without loss of generality, that  $c(x_1) = 1_a$  and  $c(x_2) = 1_b$ . We have  $1_b \in C_1(x_1)$  and  $1_a \in C_1(x_2)$  (if not we would be in Case 1 by recoloring  $x_1$  or  $x_2$ ). Recolor every vertex at distance at most 3 from x colored 3 by its subsidiary color (by Property iii), such a vertex has no sibling colored 2) and set c(x) = 3.
- Case 5.  $C_1(x) = \{1_a, \alpha\}$  and  $C_1(y) = \{1_a, \beta\}$ , with  $\alpha \in \{2_a, 2_b, 3\}$ ,  $\beta \in \{2_a, 2_b, 3\}$ . Then set  $c(y) = 1_b$ . Suppose  $c(x_1) = 1_a$  and  $c(x_2) = \alpha$ ,  $x_1$  has a neighbor colored  $1_b$  (if not we would be in Case by recoloring  $x_1$ ) and  $\{1_a, 1_b\} \in C_1(x_2)$ , by Property i). If  $\alpha = 3$ , then change the color of  $x_2$  by its subsidiary color  $\alpha'$ . Let  $\gamma \in \{2_a, 2_b\}$ , with  $\gamma \neq \alpha$ , and  $\gamma \neq \alpha'$  if  $\alpha = 3$ . Recolor every vertex at distance at most 2 from x colored  $\gamma$  by its subsidiary color and set  $c(x) = \gamma$ .
- Case 6.  $C_1(x) = \{1_a, 1_b\}$  and  $C_1(y) = \{1_a, \alpha\}$ , with  $\alpha \in \{2_a, 2_b, 3\}$ . Then set  $c(y) = 1_b$ . Suppose  $c(x_1) = 1_a$  and  $c(x_2) = 1_b$ ,  $x_1$  has a neighbor colored  $1_b$  (if not we would be in Case 1 by recoloring  $x_1$ ) and  $x_2$  has neighbors colored  $1_a$  (if not we would be in Case 2 by recoloring  $x_2$ ). Recolor every vertex at distance at most 2 from x colored  $x_2$  by its subsidiary color and set  $x_2$ .
- Case 7.  $C_1(x) = \{1_a, 1_b\}$  and  $C_1(y) = \{1_a, 1_b\}$ . Suppose  $c(x_1) = 1_a$  and  $c(x_2) = 1_b$ ,  $x_1$  has a neighbor colored  $1_b$  and  $x_2$  has a neighbor  $1_a$  (if not we would be in Case 3 by recoloring  $x_1$  or  $x_2$ ). Recolor each neighbor of  $x_1$  or  $x_2$  colored  $x_2$  by its subsidiary color (its sibling is colored 1) and set  $x_1$  and set  $x_2$  colored  $x_2$  by its subsidiary color (its sibling is colored 1) and set  $x_2$  colored  $x_2$  by its subsidiary color (its sibling is colored 1) and set  $x_2$  colored  $x_2$  by its subsidiary color (its sibling is colored 1) and set  $x_2$  colored  $x_2$  by its subsidiary color (its sibling is colored 1) and set  $x_2$  colored  $x_2$  by its subsidiary color (its sibling is colored 1) and set  $x_2$  colored  $x_2$  by its subsidiary color (its sibling is colored 1) and set  $x_2$  colored 1.

Therefore, in each case, we obtain a (1,1,2,2,3)-coloring of G.

The Petersen graph is an example of cubic graph which is not (1,1,2,3)-colorable, showing that the result of Proposition 6 is tight in a certain sense. However, experiments suggest that the Petersen graph could be the only non (1,1,2,3)-colorable subcubic graph, see Table 2.

Furthermore, the next result shows that the two colors 2 cannot be replaced by two colors 3 in the previous proposition.

**Proposition 7.** There exist cubic graphs different from the Petersen graph that are not (1, 1, 3, 3, 3)-colorable.

*Proof.* Consider the cubic graph depicted in Figure 7. Since it has diameter 3, hence no more than one vertex could be colored by a color 3. Moreover, it contains four triangles

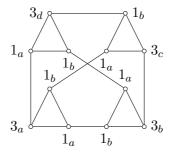


Figure 7: A cubic non (1, 1, 3, 3, 3)-colorable graph of order 12.

and each triangle should contain one vertex not colored 1. Thus, it is impossible to color it with the sequence (1, 1, 3, 3, 3).

$n \backslash S$	(1, 1)	(1, 1, 2)	(1, 1, 2, 3)	(1, 1, 2, 3, 3)
4	0	0	1	0
6	1	0	1	0
8	1	2	2	0
10	2	9	7	1
12	5	42	38	0
14	13	314	182	0
16	38	2808	1214	0
18	149	32766	8386	0
20	703	423338	86448	0
22	4132	6212201	1103114	0

Table 2: Number of S-chromatic cubic graphs of order n up to 22.

We now show that 3-irregular subcubic graphs are (1,1,2)-colorable. For subdivided graphs S(G) of any graph G, note that S(G) is (1,1)-colorable as it is bipartite.

**Proposition 8.** Every 3-irregular subcubic graph is (1,1,2)-colorable.

*Proof.* Let G be a 3-irregular graph and let e = xy be any edge of G such that x and y both have degree at most 2. If no such edge exist then the graph is bipartite and consequently (1,1)-colorable.

Define a level ordering  $L_i$ ,  $0 \le i \le r = \epsilon(e)$ , of (G, e).

We first construct a coloring c of the vertices of G from level r to 1 and with colors from the set  $\{1_a, 1_b, 2\}$ , that satisfies the following property:

i) No vertex of degree at most 2 is colored 2.

The set  $L_r$  induces a disjoint union of paths of order at most 3 in G. Since paths are (1,1)-colorable,  $L_r$  is (1,1)-colorable. Thus, Property i) is satisfied.

Assume that we have already colored all vertices of G of levels from r to i+1 and that we are going to color vertex  $u \in L_i$ ,  $1 \le i \le r-1$ . If u has degree at most 2 then  $\{1_a, 1_b\} \not\subseteq C_1(u)$ . Hence u can be colored  $1_a$  or  $1_b$  and Property i) is satisfied. If u has degree 3 and if  $C_1(u) \ne \{1_a, 1_b\}$ , then u can be colored  $1_a$  or  $1_b$ . Else if  $C_1(u) = \{1_a, 1_b\}$ , let  $u_1$  and  $u_2$  be the colored neighbors of u, with  $c(u_1) = 1_a$  and  $c(u_2) = 1_b$ . The vertex  $u_1$  has a neighbor colored  $1_b$  and the vertex  $u_2$  has a neighbor colored  $1_a$ , if not  $u_1$  and  $u_2$  could be recolored and u could be colored  $1_a$  or  $1_b$ . Thus, we have  $C_1(u) \cup C_2(u) = \{1_a, 1_b\}$  and we can color u by the color 2.

Finally, it remains to color vertices of  $L_0$ , i.e. x and y. If  $1_a \notin C_1(x)$  and  $1_b \notin C_1(y)$  (or, symmetrically,  $1_b \notin C_1(x)$  and  $1_a \notin C_1(y)$ ). Then set  $c(x) = 1_a$  and  $c(y) = 1_b$  (or, symmetrically,  $c(x) = 1_b$  and  $c(y) = 1_a$ ). Let  $x_1$  be the possible neighbor of x different from y and let  $y_1$  be the possible neighbor of y different from x. Without loss of generality, suppose that  $c(x_1) = 1_a$  and  $c(y_1) = 1_a$ . Suppose that  $x_1$  has degree at most 2. If  $C_1(x_1) = \{2\}$ , then  $x_1$  can be recolored  $1_b$  and we can set  $c(x) = 1_a$  and  $c(y) = 1_b$ . Else,  $C_1(x_1) = \{1_b\}$  by Property i) and we can set c(x) = 2 and  $c(y) = 1_b$ . If  $x_1$  has degree 3, then every colored neighbor of x has at most degree 2 and is colored  $1_b$  by Property i). Thus, as  $2 \notin C_1(x_1)$ , we can set c(x) = 2 and  $c(y) = 1_b$ . Therefore, we obtain a (1, 1, 2)-coloring of G.

## 4 (1, 2, 3, ...)-coloring

The question of whether cubics graphs have finite packing chromatic number or not was raised by Goddard et al. [11]. We give some partial results related to this question.

For the subdivision of a cubic graph, Proposition 1 along with Proposition 6 allow to obtain the following corollaries:

Corollary 3. For every subcubic graph G, S(G) is (1,3,3,5,5,7)-colorable.

Corollary 4. For every subcubic graph G,  $\chi_{\rho}(S(G)) \leq 6$ .

On the other side, it can be easily verified that  $\chi_{\rho}(S(K_4)) = 5$ . For arbitrary cubic graphs, we can (only) state the following:

**Proposition 9.** There exists a cubic graph with packing chromatic number 13.

*Proof.* The cubic graph of order 38 and diameter 4 (which is the largest cubic graph with diameter 4) described independently in [1, 18] needs 13 colors to be packing colored (checked by computer). By running a brute force search algorithm, we found that at most 28 vertices can be colored with colors  $\{1, 2, 3\}$ . But, since this graph has diameter 4, then every color greater than 3 can be given to only one vertex, implying the use of all colors from  $\{4, \ldots, 13\}$  to complete the coloring.

The distribution of packing chromatic numbers for cubic graphs of order up to 20 is presented in Table 3. We also found, (with the help of a computer), a cubic graph of order 24 and packing chromatic number 11.

$n\setminus\chi_{ ho}$	4	5	6	7	8	9	10	11
4	1	0	0	0	0	0	0	0
6	1	1	0	0	0	0	0	0
8	0	3	2	0	0	0	0	0
10	0	3	15	1	0	0	0	0
12	0	7	42	36	0	0	0	0
14	0	13	252	222	22	0	0	0
16	0	34	907	2685	433	1	0	0
18	0	116	5277	21544	14050	314	0	0
20	0	151	22098	206334	226622	5528	84*	0

Table 3: Number of cubic graphs of order n with packing chromatic number  $\chi_{\rho}$  up to 20.\*There are 55284 cubic graphs of order 20 and with packing chromatic number between 9 and 10 (our program takes too long time to compute their packing chromatic numbers).

## 5 Concluding remarks

We conclude this paper by listing a few open problems:

- Is it true that any subcubic graph except the Petersen graph is (1, 1, 2, 3)-colorable?
- Is it true that any subcubic graph except the Petersen graph is (1, 2, 2, 2, 2, 2)colorable?
- Does there exist a 3-irregular subcubic graph that is not (1, 2, 2, 3)-colorable?
- Is it true that any 3-irregular subcubic graph is (1, 1, 3)-colorable?
- Is it true that the subdivision of any subcubic graph is (1, 2, 3, 4, 5)-colorable?
- Does there exist a cubic graph with packing chromatic number larger than 13?

#### References

- [1] I. Alegre, M. A. Fiol, and J. L. A. Yebra. Some large graphs with given degree and diameter. *J. Graph Theory*, 10(2):219–224, 1986.
- [2] Oleg V. Borodin and Anna O. Ivanova. 2-Distance 4-coloring of planar subcubic graphs J. Appl. Ind. Math., 5(4):535–541, 2011.
- [3] Boštjan Brešar, Sandi Klavžar, and Douglas F. Rall. On the packing chromatic number of Cartesian products, hexagonal lattice, and trees. *Discrete Appl. Math.*, 155(17):2303–2311, 2007.
- [4] Daniel W. Cranston and Seog-Jin Kim. List-coloring the square of a subcubic graph. J. Graph Theory, 57(1):65–87, 2008.

- [5] Jan Ekstein, Jiří Fiala, Premysl Holub, and Bernard Lidický. The packing chromatic number of the square lattice is at least 12. arXiv:1003.2291v1, 2010.
- [6] Jiří Fiala and Petr A. Golovach. Complexity of the packing coloring problem for trees. *Discrete Appl. Math.*, 158(7):771–778, 2010.
- [7] Jiří Fiala, Sandi Klavžar, and Bernard Lidický. The packing chromatic number of infinite product graphs. *European J. Combin.*, 30(5):1101–1113, 2009.
- [8] Arthur S. Finbow and Douglas F. Rall. On the packing chromatic number of some lattices. *Discrete Appl. Math.*, 158(12):1224–1228, 2010.
- [9] Nicolas Gastineau. On dichotomies among the instances of the S-coloring problem. Manuscript, 2013.
- [10] Nicolas Gastineau, Hamamache Kheddouci, and Olivier Togni. Subdivision into *i*-packings and S-packing chromatic number of some lattices. Manuscript, 2013.
- [11] Wayne Goddard, Sandra M. Hedetniemi, Stephen T. Hedetniemi, John M. Harris, and Douglas F. Rall. Broadcast chromatic numbers of graphs. Ars Combin., 86:33–49, 2008.
- [12] Wayne Goddard and Honghai Xu. A note on packing colorings of lattices. *Discrete Appl. Math.*, 166:255–262, 2014.
- [13] Wayne Goddard and Honghai Xu. The S-packing chromatic number of a graph. Discuss. Math. Graph Theory, 32(4):795–806, 2012.
- [14] Gordon Royle. List of cubic graphs. http://mapleta.maths.uwa.edu.au/~gordon/remote/cubics/, last accessed March 26, 2014.
- [15] Frédéric Havet Choosability of the square of planar subcubic graphs with large girth. *Discrete Math.*, 309(11):3553-3563, 2009.
- [16] Florica Kramer and Horst Kramer. A survey on the distance-colouring of graphs. *Discrete Math.*, 308(2–3):422–426, 2008.
- [17] Roman Soukal and Premysl Holub. A note on packing chromatic number of the square lattice. *Electron. J. Combin.*, (N17), 2010.
- [18] C. Von Conta. Torus and other networks as communication networks with up to some hundred points. *IEEE Trans. Comput.*, C-32(7):657 666, 1983.